

Figure 1: Schematic representation of the experimental design. The figure shows a flowchart of the experimental design. It starts with 'Pretest' (N=10) leading to 'Main Experiment' (N=100). The 'Main Experiment' is divided into 'Control' (N=50) and 'Treatment' (N=50). The 'Control' group is further divided into 'Control 1' (N=25) and 'Control 2' (N=25). The 'Treatment' group is divided into 'Treatment 1' (N=25) and 'Treatment 2' (N=25). The 'Treatment 1' group is further divided into 'Treatment 1a' (N=12.5) and 'Treatment 1b' (N=12.5). The 'Treatment 2' group is further divided into 'Treatment 2a' (N=12.5) and 'Treatment 2b' (N=12.5). The 'Treatment 1a' and 'Treatment 2a' groups are further divided into 'Treatment 1a1' (N=6.25) and 'Treatment 1a2' (N=6.25), and 'Treatment 2a1' (N=6.25) and 'Treatment 2a2' (N=6.25) respectively. The 'Treatment 1b' and 'Treatment 2b' groups are further divided into 'Treatment 1b1' (N=6.25) and 'Treatment 1b2' (N=6.25), and 'Treatment 2b1' (N=6.25) and 'Treatment 2b2' (N=6.25) respectively. The 'Treatment 1a1' and 'Treatment 1a2' groups are further divided into 'Treatment 1a1a' (N=3.125) and 'Treatment 1a1b' (N=3.125), and 'Treatment 1a2a' (N=3.125) and 'Treatment 1a2b' (N=3.125) respectively. The 'Treatment 1b1' and 'Treatment 1b2' groups are further divided into 'Treatment 1b1a' (N=3.125) and 'Treatment 1b1b' (N=3.125), and 'Treatment 1b2a' (N=3.125) and 'Treatment 1b2b' (N=3.125) respectively. The 'Treatment 2a1' and 'Treatment 2a2' groups are further divided into 'Treatment 2a1a' (N=3.125) and 'Treatment 2a1b' (N=3.125), and 'Treatment 2a2a' (N=3.125) and 'Treatment 2a2b' (N=3.125) respectively. The 'Treatment 2b1' and 'Treatment 2b2' groups are further divided into 'Treatment 2b1a' (N=3.125) and 'Treatment 2b1b' (N=3.125), and 'Treatment 2b2a' (N=3.125) and 'Treatment 2b2b' (N=3.125) respectively. The 'Treatment 1a1a' and 'Treatment 1a1b' groups are further divided into 'Treatment 1a1a1' (N=1.5625) and 'Treatment 1a1a2' (N=1.5625), and 'Treatment 1a1b1' (N=1.5625) and 'Treatment 1a1b2' (N=1.5625) respectively. The 'Treatment 1a2a' and 'Treatment 1a2b' groups are further divided into 'Treatment 1a2a1' (N=1.5625) and 'Treatment 1a2a2' (N=1.5625), and 'Treatment 1a2b1' (N=1.5625) and 'Treatment 1a2b2' (N=1.5625) respectively. The 'Treatment 1b1a' and 'Treatment 1b1b' groups are further divided into 'Treatment 1b1a1' (N=1.5625) and 'Treatment 1b1a2' (N=1.5625), and 'Treatment 1b1b1' (N=1.5625) and 'Treatment 1b1b2' (N=1.5625) respectively. The 'Treatment 1b2a' and 'Treatment 1b2b' groups are further divided into 'Treatment 1b2a1' (N=1.5625) and 'Treatment 1b2a2' (N=1.5625), and 'Treatment 1b2b1' (N=1.5625) and 'Treatment 1b2b2' (N=1.5625) respectively. The 'Treatment 2a1a' and 'Treatment 2a1b' groups are further divided into 'Treatment 2a1a1' (N=1.5625) and 'Treatment 2a1a2' (N=1.5625), and 'Treatment 2a1b1' (N=1.5625) and 'Treatment 2a1b2' (N=1.5625) respectively. The 'Treatment 2a2a' and 'Treatment 2a2b' groups are further divided into 'Treatment 2a2a1' (N=1.5625) and 'Treatment 2a2a2' (N=1.5625), and 'Treatment 2a2b1' (N=1.5625) and 'Treatment 2a2b2' (N=1.5625) respectively. The 'Treatment 2b1a' and 'Treatment 2b1b' groups are further divided into 'Treatment 2b1a1' (N=1.5625) and 'Treatment 2b1a2' (N=1.5625), and 'Treatment 2b1b1' (N=1.5625) and 'Treatment 2b1b2' (N=1.5625) respectively. The 'Treatment 2b2a' and 'Treatment 2b2b' groups are further divided into 'Treatment 2b2a1' (N=1.5625) and 'Treatment 2b2a2' (N=1.5625), and 'Treatment 2b2b1' (N=1.5625) and 'Treatment 2b2b2' (N=1.5625) respectively.

**CROSS REFERENCE TO RELATED APPLICATIONS**

The present application is related to co-pending U.S. Patent Application Serial No. \_\_\_\_\_ (IBM Docket No. AUS920010398US1) entitled "Method and Apparatus for Wide-Spread Distribution of Electronic Content in a Non-Linear Peer to Peer Fashion" filed even date herewith. The content of the above mentioned commonly assigned, co-pending U. S. Patent applications are hereby incorporated herein by reference for all purposes.

**1. Technical Field:**

## 2. Description of Related Art:

Current technology for mass distribution of data over the Internet consists of one or more "master" servers where the content is available, and many more "mirror" sites where the same data is stored. These mirror sites do not actually use the data themselves. Typically, the master server is overwhelmed very easily, and end users are forced to manually attempt a list of mirror sites. Each of those mirror sites may or may not actually have the updated content because they are typically driven by time-based automation (typically a

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cron job scheduled at midnight). This distribution scheme is incredibly problematic and wasteful in dealing with the initial wave of interest in specific data.

Many of these problems may be reduced by using  
5 peer-to-peer technology to offload demands from master servers to other nearby clients which are downloading the same content for their own use. The master server divides a large file into several small pieces and then downloads those file pieces to the first client machines  
10 which request the file. For example, a 50 megabyte (MB) file may be broken into 50 1-MB pieces which are then downloaded to 50 different clients. These clients will then function as peer-to-peer servers. Subsequent requests from new client machines are then redirected by  
15 the master server to the clients which already have the required file pieces.

Because the client machine are owned by the end users, acquiring peer-to-peer connections may be a significant problem if people do not want to share their  
20 bandwidth and computer resources.

Therefore, it would be desirable to have a method for providing incentive for end users to allow their client machine to act as peer-to-peer servers.

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**SUMMARY OF THE INVENTION**

The present invention provides a method, program and system to provide incentives for client machines to contribute resources to a peer-to-peer computer network. When a server receives requests for information from a plurality of client machines, it determines if the client machines are contributing resources to peer-to-peer sharing. When answering requests, clients which contribute resources to peer-to-peer sharing are given priority over clients which do not contribute. In another embodiment of the present invention, a further incentive is provided to clients which contribute to peer-to-peer sharing, by giving higher priority to client requests in proportion to the level of resources contributed.

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The novel features believed characteristic of the invention are set forth in the appended claims. The invention itself, however, as well as a preferred mode of use, further objectives and advantages thereof, will best be understood by reference to the following detailed description of an illustrative embodiment when read in conjunction with the accompanying drawings, wherein:

**Figure 2** depicts a block diagram of a data processing system that may be implemented as a server in accordance with a preferred embodiment of the present invention;

20           **Figure 4** depicts a flowchart illustrating a method  
for encouraging clients to contribute resources to  
peer-to-peer sharing in accordance with the present  
invention.

**DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT**

With reference now to the figures, **Figure 1** depicts a pictorial representation of a network of data processing systems in which the present invention may be implemented. Network data processing system **100** is a network of computers in which the present invention may be implemented. Network data processing system **100** contains a network **102**, which is the medium used to provide communications links between various devices and computers connected together within network data processing system **100**. Network **102** may include connections, such as wire, wireless communication links, or fiber optic cables.

In the depicted example, a server **104** is connected to network **102** along with storage unit **106**. In addition, clients **108**, **110**, and **112** also are connected to network **102**. These clients **108**, **110**, and **112** may be, for example, personal computers or network computers. In the depicted example, server **104** provides data, such as boot files, operating system images, and applications to clients **108-112**. Clients **108**, **110**, and **112** are clients to server **104**. Network data processing system **100** may include additional servers, clients, and other devices not shown.

In the depicted example, network data processing system **100** is the Internet with network **102** representing a worldwide collection of networks and gateways that use the TCP/IP suite of protocols to communicate with one another. At the heart of the Internet is a backbone of high-speed data communication lines between major nodes or host computers, consisting of thousands of commercial, government, educational and other computer systems that

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route data and messages. Of course, network data processing system **100** also may be implemented as a number of different types of networks, such as for example, an intranet, a local area network (LAN), or a wide area network (WAN). **Figure 1** is intended as an example, and not as an architectural limitation for the present invention.

Referring to **Figure 2**, a block diagram of a data processing system that may be implemented as a server, such as server **104** in **Figure 1**, is depicted in accordance with a preferred embodiment of the present invention. Data processing system **200** may be a symmetric multiprocessor (SMP) system including a plurality of processors **202** and **204** connected to system bus **206**. Alternatively, a single processor system may be employed. Also connected to system bus **206** is memory controller/cache **208**, which provides an interface to local memory **209**. I/O bus bridge **210** is connected to system bus **206** and provides an interface to I/O bus **212**. Memory controller/cache **208** and I/O bus bridge **210** may be integrated as depicted.

Peripheral component interconnect (PCI) bus bridge **214** connected to I/O bus **212** provides an interface to PCI local bus **216**. A number of modems may be connected to PCI bus **216**. Typical PCI bus implementations will support four PCI expansion slots or add-in connectors. Communications links to network computers **108-112** in **Figure 1** may be provided through modem **218** and network adapter **220** connected to PCI local bus **216** through add-in boards.

Additional PCI bus bridges **222** and **224** provide interfaces for additional PCI buses **226** and **228**, from which additional modems or network adapters may be supported. In this manner, data processing system **200** allows connections to multiple network computers. A memory-mapped graphics adapter **230** and hard disk **232** may also be connected to I/O bus **212** as depicted, either directly or indirectly.

Those of ordinary skill in the art will appreciate that the hardware depicted in **Figure 2** may vary. For example, other peripheral devices, such as optical disk drives and the like, also may be used in addition to or in place of the hardware depicted. The depicted example is not meant to imply architectural limitations with respect to the present invention.

The data processing system depicted in **Figure 2** may be, for example, an IBM RISC/System 6000 system, a product of International Business Machines Corporation in Armonk, New York, running the Advanced Interactive Executive (AIX) operating system.

With reference now to **Figure 3**, a block diagram illustrating a data processing system is depicted in which the present invention may be implemented. Data processing system **300** is an example of a client computer. Data processing system **300** employs a peripheral component interconnect (PCI) local bus architecture. Although the depicted example employs a PCI bus, other bus architectures such as Accelerated Graphics Port (AGP) and Industry Standard Architecture (ISA) may be used. Processor **302** and main memory **304** are connected to PCI local bus **306** through PCI bridge **308**. PCI bridge **308** also may include an integrated memory controller and cache

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memory for processor **302**. Additional connections to PCI local bus **306** may be made through direct component interconnection or through add-in boards. In the depicted example, local area network (LAN) adapter **310**, SCSI host bus adapter **312**, and expansion bus interface **314** are connected to PCI local bus **306** by direct component connection. In contrast, audio adapter **316**, graphics adapter **318**, and audio/video adapter **319** are connected to PCI local bus **306** by add-in boards inserted into expansion slots. Expansion bus interface **314** provides a connection for a keyboard and mouse adapter **320**, modem **322**, and additional memory **324**. Small computer system interface (SCSI) host bus adapter **312** provides a connection for hard disk drive **326**, tape drive **328**, CD-ROM drive **330**, and DVD drive **332**. Typical PCI local bus implementations will support three or four PCI expansion slots or add-in connectors.

An operating system runs on processor **302** and is used to coordinate and provide control of various components within data processing system **300** in **Figure 3**. The operating system may be a commercially available operating system, such as Windows 2000, which is available from Microsoft Corporation. An object oriented programming system such as Java may run in conjunction with the operating system and provide calls to the operating system from Java programs or applications executing on data processing system **300**. "Java" is a trademark of Sun Microsystems, Inc. Instructions for the operating system, the object-oriented operating system, and applications or



programs are located on storage devices, such as hard disk drive **326**, and may be loaded into main memory **304** for execution by processor **302**.

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incentives for the owners of client machines to contribute their resources to peer-to-peer sharing technology.

Referring to **Figure 4**, a flowchart illustrating a method for encouraging clients to contribute resources to peer-to-peer sharing is depicted in accordance with the present invention. On the master server, requests for files will be answered first for peer-to-peer servers and secondly for clients in the normal fashion used today.

When the master server receives a file request from a client (step **401**), the master determines if that client has adopted peer-to-peer sharing technology (step **402**). If the client is not contributing resources for peer-to-peer sharing, the file request will still be granted, but the request is placed into the "slow lane" of the pending file transfers (step **403**). If the client is contributing resources for peer-to-peer sharing, the master server places the request in the "fast lane" (step **404**). Client requests given fast lane status are always given higher priority than requests with slow lane status.

In contributing to peer-to-peer sharing, clients would be allowed to delegate a "sandbox" for the protocol to use. This includes specific disk space limits, bandwidth limits, CPU limits, memory limits, and limits on number of users connecting.

The caveat of the present approach is that some users may select minimal resources necessary in order to get into the fast lane of the peer-to-peer servers. To address the problem, the master server maintains a priority queue for requests that are in the fast lane. This priority queue evaluates the level of resources that



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The description of the present invention has been presented for purposes of illustration and description, and is not intended to be exhaustive or limited to the invention in the form disclosed. Many modifications and variations will be apparent to those of ordinary skill in the art. The embodiment was chosen and described in order to best explain the principles of the invention, the practical application, and to enable others of ordinary skill in the art to understand the invention for various embodiments with various modifications as are suited to the particular use contemplated.

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